



## GVDS Assets

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- ◆ In coordination models exploiting the notion of GVDS:
  - ◆ The association between the coordination context contributed by a given agent and the agent itself is now made explicit
  - ◆ The resulting style of coordination draws a distinction between the information immediately available to an agent and the one that can be requested from others
  - ◆ Still, the benefits of coordination, e.g., the decoupling of communication from behavior, are retained
- ◆ Hence, GVDS fosters a coordination style where:
  - ◆ coordination is defined entirely in terms of the coordinated agents, without reliance on some external entity
  - ◆ the coordination context is automatically and dynamically reconfigured
  - ◆ coordination is achieved through local actions that have a global effect
- ◆ The conjecture is that these characteristics are going to:
  - ◆ simplify the task of building (and reasoning about) applications that ...
  - ◆ ... are built out of autonomous components ...
  - ◆ ... whose relationships are dynamically and frequently reconfigured

2



## Design Alternatives

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- ◆ Choice of the data structure
  - ◆ Sets, bags, trees, graphs, matrices, ...
  - ◆ May affect the efficiency and/or complexity of the implementation
- ◆ Choice of operations
  - ◆ Local vs. global
  - ◆ Query vs. manipulation
  - ◆ Proactive vs. reactive
  - ◆ Synchronous vs. asynchronous
- ◆ Choice of the partitioning/merging criteria
  - ◆ Superposition, union, composition, ...
- ◆ Choice of the enabling condition for sharing
  - ◆ Based on connectivity
    - connectivity over space vs. connectivity over time for physical mobility
    - co-location for logical mobility
  - ◆ Possibly augmented by application constraints
    - e.g., to deal with security, or with specific application constraints

3



## Design Alternatives – cont'd

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- ◆ Degree of symmetry and transitivity
  - ◆ Is everybody “seeing” the same content?
- ◆ Degree of atomicity
  - ◆ Important when determining the semantics of operations, and their relationship to sharing
  - ◆ Determines the extent to which one can treat the GVDS as a “local” data structure
  - ◆ Simplifying the programmer’s chore vs. delivering an efficient implementation
- ◆ Degree of consistency
  - ◆ Given two agents, how far can their perception of the GVDS drift?
  - ◆ The answer to this question often implies the use of caching and replication schemes
- ◆ Degree of knowledge about the system configuration
  - ◆ System information can be represented in a GVDS, too
- ◆ Degree of persistency
  - ◆ If a portion of the system is known to be stable, how can we exploit it?

4



## Research Issues

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- ◆ What is the good balance to strike among the design alternatives?
  - ◆ Relationship with other middleware approaches and results
- ◆ Is there a “unifying theory” of GVDS?
  - ◆ Is it possible to separate the issues related with distribution from those intimately connected to the data structure chosen?
  - ◆ A positive answer could lead to a middleware supporting instantiations of GVDS with different data structures
- ◆ What is the relationship between GVDS and security?
- ◆ What is the impact of the GVDS abstraction on formal reasoning and verification?

5



## GVDS Summary

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- ◆ Global virtual data structures are a novel coordination paradigm targeted at highly dynamic environments
- ◆ GVDS is not meant to be a new model by itself: instead, it is meant to be the driving concept behind a new family of coordination models
- ◆ While some incarnations of GVDS are already available, only a fraction of the design space has been explored so far

6



## Outline

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- ◆ Introduction and Major Issues
- ◆ Commercial Mobile Middleware
- ◆ Next-Generation Mobile Middleware
- ◆ Case Study – LIME
- ◆ Middleware for Wireless Sensor Networks
- ◆ Open Issues and Future Directions

7



# Service Discovery

- ◆ When enter a new area, problem is how to identify and connect with services
- ◆ Part of adaptability, and part of almost every mobile middleware system
  - ◆ e.g., replication for ad hoc collaborative environments used SD for identifying other devices.
  - ◆ e.g., Spectra (part of Aura) uses service discovery to find other computational devices to spread computation
- ◆ Goal: not to re-implement service discovery for every system (including middleware), but re-use a general solution
  - ◆ SLP, UPnP, SDP (Bluetooth), Jini

8



# Service Location Protocol (SLP)

IETF

Service Discovery

- ◆ Standards driven – developed within IETF
- ◆ Decentralized, lightweight, scalable and extensible protocol for service discovery within a site.
- ◆ Directory agent (optional), Service agent, User agent
- ◆ Active and passive DA discovery
  - ◆ DHCP option 78: distribution of DA addresses
- ◆ Services identified by URL
  - ◆ Service templates – exploit standardized vocabulary
  - ◆ Connection mechanism independent from SLP



9



# Universal Plug and Play (UPnP)

Microsoft

Service Discovery

- ◆ Microsoft lead industry consortium (617 members as of September 2003)
- ◆ Discovery of devices reachable through TCI/IP
- ◆ No central service registry
  - ◆ Control points (e.g., your Windows ME laptop) "search" for devices (multicast query)
  - ◆ Control achieved through web page or specialized interface (e.g., control your new DVD player through its web page, rather than the remote control)
- ◆ Devices *multicast* their services
- ◆ XML based service descriptions

10



# Service Discovery Protocol (SDP)

Bluetooth

Service Discovery

- ◆ Focuses on service discovery for ad hoc networks
  - ◆ No central service registry
- ◆ Three search types (on existing connection!):
  - ◆ By service type
  - ◆ By service attribute
  - ◆ Or service browsing
- ◆ Can retrieve attributes that detail how to connect to the service (protocol identifier, and protocol-specific parameters)
- ◆ Does not "compete" with other discovery protocols, but can have them layered on top to provide additional functionality (e.g., leasing, security, etc.)

11



# Jini Service Discovery

Sun

Service Discovery

- ◆ Sun Microsystems, open source
- ◆ Mechanisms for discovery of lookup server (required component)
  - ◆ *Multicast request protocol*
    - Browse for local, nearby lookup services, UDP every 5 sec.
  - ◆ *Multicast announcement protocol*
    - Lookup services advertise existence, clients listen
    - Used by new lookup services or after a failure recovery
  - ◆ *Unicast discovery protocol*
    - Join a specific community, either known, or discovered by above
- ◆ Result of service discovery is Jini object code for direct access to the server

12



# SD Summary

- ◆ Many competing industrial standards
- ◆ Tradeoff between centralized registry and “noise” of protocol
  - ◆ Multicast increases network traffic, which is a concern for mobile, battery-powered devices
- ◆ Key component often missing is the vocabulary used for service description
  - ◆ XML is extensible, but still need shared vocabulary in order to work together

13



## Event-Based

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- ◆ Publish-Subscribe systems are asynchronous, implicit, multi-point, and peer-to-peer in communication style
- ◆ This style is suited to both traditional distributed systems and to mobile systems
- ◆ Clients and publishers are decoupled, and the infrastructure can be distributed

14



## Solar

Dartmouth College

Event-Based

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- ◆ Context-information collection, aggregation, and dissemination (data fusion)
  - ◆ Operator graph representation of computation allows decomposition and reuse of context aggregation primitives
  - ◆ Examples: filters, transformers (lookup mechanisms)
- ◆ Applications register operations with centralized server "Star" (the centralized dispatcher)
- ◆ Computation farmed to available hosts (planets)
  - ◆ Directly deliver events to applications

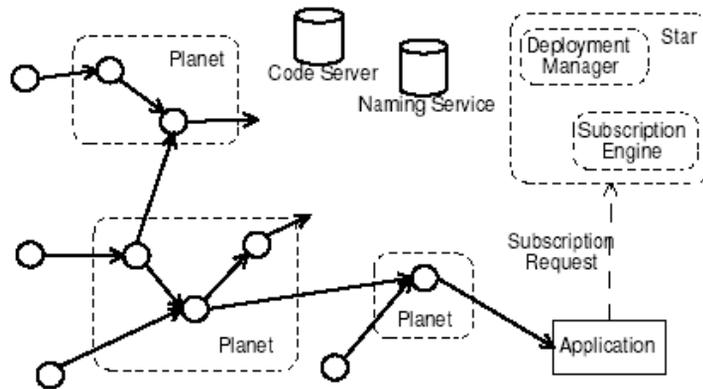
15



# Solar Architecture

Dartmouth College

Event-Based



16



# Jedi

Polimi

Event-Based

- ◆ Provides a scalable, distributed, content-based event dispatcher
- ◆ Supports mobile agents that connect to a dispatcher
  - ◆ Clients disconnect
  - ◆ The old event dispatcher stores events for disconnected clients
  - ◆ When client reconnects, stored events are transferred and delivered
  - ◆ Partial order of events is guaranteed

17



# Reconfiguration of PS Infrastructure

- ◆ A distributed dispatcher is potentially useful in a mobile ad hoc environment.
  - ◆ Disconnected nature of PS matches requirements and open nature of MANET
  - ◆ Each host becomes a node in the dispatching tree
- ◆ Current solutions do not address reconfiguration of the infrastructure itself
  - ◆ Reconfiguration is required when nodes fail, when links disappear, when the PS tree is partitioned
- ◆ We are examining solutions that adapt an MAODV-like tree reconstruction algorithm for fixing the PS infrastructure in MANET environments



# PS Summary

- ◆ Although pub/sub features match the mobile environment well, the challenging issues have not all been addressed
- ◆ There seem to be strong connections between the tree of pub/sub and the tree of multicast MANET solutions
  - ◆ We are exploring this connection, and hope that we can develop pub/sub successfully in a MANET



# Object-Oriented

- ◆ Communication among distributed objects
- ◆ Applicability of existing OO middleware to mobile setting is limited
  - ◆ Heavy computational load (need to be made lighter!)
  - ◆ Synchronous communication
  - ◆ Transparency (might be difficult to be light and adaptive without optimizing considering application needs)
- ◆ Risk to penalize performance, generate high costs and lead to unusable and non-scalable solutions

20



# Rover

MIT

Object Oriented

- ◆ Aims to support both transparency and awareness, and is based on two key concepts:
  - ◆ **Relocatable dynamic objects (RDO)**: mobile objects (code and data) that can be dynamically loaded into a client from a server, or vice versa
    - Essentially used to reduce communication, e.g., to perform semantic compression
    - Methods contain also conflict resolution procedures
  - ◆ **Queued remote procedure calls (QRPC)**: provides the core communication mechanisms among RDOs, by providing a non-blocking RPCs
- ◆ A client-server architecture is assumed, no mobile-mobile interaction supported
- ◆ The toolkit has been tested on the field by porting several common applications (e.g., `exmh`, `ical`, and other X11 applications)
  - ◆ A speed-up of about 20% has been measured, but mostly for scenarios with very low bandwidth

21



# Using RDOs

MIT - Rover

Object Oriented

- ◆ A Rover application is made of RDOs, that are fetched lazily from servers
  - ◆ Each RDO has a "home" server
- ◆ RDOs are cached and replicated
  - ◆ the primary copy is held at the server
  - ◆ a prefetching scheme is employed to fill the object cache while connected with objects that are selected by the application
- ◆ Updates to local copies of RDOs are done immediately, and marked tentative
  - ◆ they are lazily propagated to servers using QRPC
  - ◆ It is up to the applications to choose to use tentative RDOs
- ◆ Servers are in charge of detecting and resolving conflicts
  - ◆ this is typically achieved through methods invoked on the primary copy, by keeping version vectors for RDOs, and by examining invocation logs
  - ◆ A wide range of consistency control techniques are provided, although primary-copy, tentative-updates are preferred
- ◆ The system itself is made of RDOs, that can be loaded on-demand according to the features of the current environment

22



# QRPC

MIT - Rover

Object Oriented

- ◆ QRPC is used to fetch RDOs, as well as to enable the communication between client and server, necessary to keep the object copies consistent
- ◆ Invocation of a QRPC on an object returns immediately, by yielding a *promise* to the object
  - ◆ Promise object can be used to proactively check whether a result has arrived, to suspend waiting for the result, or to register callbacks
- ◆ If a mobile host is disconnected between sending the request and receiving the reply, the server will periodically try to contact the mobile host and deliver the reply
  - ◆ Different communication channels can be used for request and reply
  - ◆ Support of disconnected operation
- ◆ The queued RPCs, issued by the client, are filtered by a network scheduler module, that:
  - ◆ may decide to deliver them in a non-FIFO order, to suit the application-specified priorities and needs
  - ◆ May compress the requests, and/or send them in batches

23



# ALICE

U. Dublin

Object-Oriented

- ◆ Architecture for Location-Independent CORBA Environments
- ◆ Mobile hosts act as servers AND/OR clients
  - ◆ Interoperate transparently with standard CORBA applications
  - ◆ Does not require support for Mobile IP
  - ◆ Hides broken TCP connections
  - ◆ Integrates service handoff with mobile handoff
- ◆ Uses IIOP
  - ◆ Minimum ORB protocols to exchange messages between clients and servers

24



## Object Oriented Summary

- ◆ Object oriented software develop is a common practice in distributed environments
- ◆ Supporting object relocation and interaction in the mobile environment eases application programming as long as the resulting semantics remain well-understood

25



# Transactions

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- ◆ Sequence of operations clustered together complying with ACID properties
  - ◆ Atomicity: execute completely or not at all
  - ◆ Consistency: lead from one consistent state to another
  - ◆ Isolation: isolation from concurrent transactions
  - ◆ Durability: once completed, changes cannot be undone
- ◆ Mobile Transactions:
  - ◆ Support mobile clients performing transactions on stationary servers
  - ◆ ACID may not be enforceable – high overhead

26



# Kangaroo

Southern Methodist U., Texas

Transactions

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- ◆ Support for transactions by mobile entities that may not start and end at the same location
- ◆ Support long-lived, hopping transactions
- ◆ Two modes
  - ◆ Split: failure of transaction at base station does not imply failure of complete transaction (previous hops)
  - ◆ Compensating: failure of a transaction at base station implies failure of complete transaction

27



## Transaction Summary

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- ◆ Most industrial transactions (e.g., Oracle Mobile agents, IBM Gold Rush) are satisfied with simple transactional support
  - ◆ Submitting remote transactions
  - ◆ Optimistically performing transactions, reporting conflicts
- ◆ Kangaroo transactions address a single form of transaction for mobile users
  - ◆ Do not solve all issues, nor provide a clean model to replace the ACID concepts
- ◆ Issues not addressed
  - ◆ Transactions and commitment in mobile ad hoc networks
  - ◆ What to do when the database itself is on a mobile device

28



## Transport Layer

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- ◆ TCP/IP does not perform well across wireless links
  - ◆ Packet loss on the lossy link cause excessive slow-start at the fixed host
- ◆ Solutions exist, including I-TCP (indirect TCP, from Rutgers U.), that remove TCPs end-to-end guarantee, but significantly improve performance for bulk transfers (1.5 times)
- ◆ Another significant problem is naming
  - ◆ All components, especially those in a disconnected MANET should not be assumed to have an IP address
  - ◆ They may not even have a distinct identifier

29



## Underlying Support Algorithms

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- ◆ Middleware often assumes certain existing technologies
- ◆ Often these technologies themselves are non-trivial, and deserve separate focus
- ◆ POMC – Principles of Mobile Computing
  - ◆ Theoretical workshop in 2001, 2002, and 2003
  - ◆ Guaranteed message delivery
  - ◆ Mutual exclusion algorithms
  - ◆ Group communication for mobile participants
  - ◆ K-cluster formation

30



## Reliable Message Delivery to Mobile Components

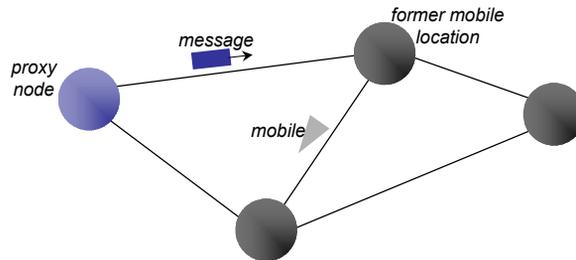
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- ◆ Applications need to communicate control and data information to mobile components
- ◆ We must have a mechanism to *reliably deliver* these messages (unicast or multicast)
- ◆ Guaranteeing the delivery of a message to a moving unit is challenging even under the assumption of a fault-free network

31



## Proxy

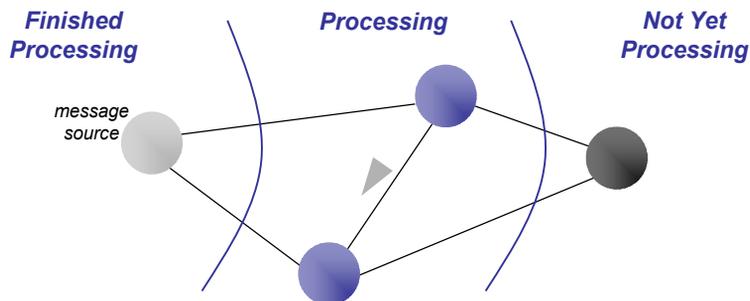


- ◆ During movement, information at the proxy is out of date and messages will chase the mobile unit
- ◆ Similar to the approach of Mobile IP

32



## Idea: Adapt Distributed Snapshot to Message Delivery



- ◆ Snapshots provide a consistent image of the network state
- ◆ Key property: ***every message in the system appears in exactly one local snapshot***

33



## From Snapshots To Message Delivery

Distributed Snapshot	Snapshot Delivery
Node	Mobile agent server/ Base station
<i>Message</i>	<i>Mobile unit</i>
Token	Application message
<i>Record message</i>	<i>Deliver application message</i>
Local snapshot terminates	Application message deleted

34



## Properties of Snapshot Delivery

- ◆ Exactly-once delivery independent of mobile movement
- ◆ Can be trivially extended to multicast
- ◆ Storage required at each node for only one message round trip time
- ◆ Overhead of one message per edge
- ◆ Network neighbors must be known in advance, but extensions allow the underlying network to be dynamic

35



## Other Algorithmic Work

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- ◆ This approach adapts distributed computing algorithms to mobile computing
  - ◆ Another solutions: using diffusing computations to track mobile units, layered with message delivery
  - ◆ Other promising algorithms: randomized algorithms, self-stabilizing algorithms, epidemic algorithms
- ◆ Other approaches include developing algorithms similar to those of distributed computing by shifting the burden of computation onto fixed infrastructure
  - ◆ Checkpointing: Acharya'94
  - ◆ Reliable multicast message delivery: Acharya'96
  - ◆ Causal event ordering: Prakash'97
  - ◆ Termination detection: Matocha'98

36



## Outline

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- ◆ Introduction and Major Issues
- ◆ Commercial Mobile Middleware
- ◆ Next-Generation Mobile Middleware
- ◆ Case Study – LIME
- ◆ Middleware for Wireless Sensor Networks
- ◆ Summary and Open Questions

37



## Summary

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- ◆ Middleware for mobile computing provides abstractions for easing the development process
- ◆ Commercial middleware is targeted toward the first step of mobility, providing service access to mobile devices
- ◆ Major issues and approaches in research include
  - ◆ Replication, Adaptation, Service Discovery, Event-based, Object-Oriented, Transactional, Transport Layer, Algorithmic, Data Sharing
  - ◆ Not discussed issues include reflection, security, ...

38



## Questions:

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- ◆ Will there be / should there be a single middleware for mobile computing?
  - ◆ Not all environments have the same demands and needs
  - ◆ Can there be a composable middleware that allows designers to pull in only the aspects that they need?
- ◆ Will the need for mobile middleware diminish as wireless networks become faster?
- ◆ Can middleware be shared among applications for:
  - ◆ Nomadic computing
  - ◆ Mobile ad hoc computing
  - ◆ Sensor networks

39